

Strong monoidal functors and modules

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I wrote this note because I wanted to know that a strong monoidal functor which is also an isomorphism of categories preserves (sub)categories of modules up to isomorphism (Proposition 11). A similar fact holds for strong monoidal equivalences (Proposition 13).

Definition 1. Let $(\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}})$ and $(\mathcal{D}, \otimes_{\mathcal{D}}, 1_{\mathcal{D}})$ be monoidal categories. A **lax monoidal functor** $(F, e, \mu_{x,y}) : (\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}}) \rightarrow (\mathcal{D}, \otimes_{\mathcal{D}}, 1_{\mathcal{D}})$ is

- a functor $F : \mathcal{C} \rightarrow \mathcal{D}$,
- a morphism $e : 1_{\mathcal{D}} \rightarrow F(1_{\mathcal{C}})$, and
- a natural transformation $\mu_{x,y} : F(x) \otimes_{\mathcal{D}} F(y) \rightarrow F(x \otimes_{\mathcal{C}} y)$,

such that

1. (**associativity**) for all $x, y, z \in \text{obj}(\mathcal{C})$, the diagram

$$\begin{array}{ccc}
 (F(x) \otimes_{\mathcal{D}} F(y)) \otimes_{\mathcal{D}} F(z) & \xrightarrow{a_{\mathcal{D}}} & F(x) \otimes_{\mathcal{D}} (F(y) \otimes_{\mathcal{D}} F(z)) \\
 \mu \otimes \text{id} \downarrow & & \downarrow \text{id} \otimes \mu \\
 F(x \otimes_{\mathcal{C}} y) \otimes_{\mathcal{D}} F(z) & & F(x) \otimes_{\mathcal{D}} F(y \otimes_{\mathcal{C}} z) \\
 \mu \downarrow & & \downarrow \mu \\
 F((x \otimes_{\mathcal{C}} y) \otimes_{\mathcal{C}} z) & \xrightarrow{F(a_{\mathcal{C}})} & F(x \otimes_{\mathcal{C}} (y \otimes_{\mathcal{C}} z))
 \end{array}$$

commutes, and

2. (**unitality**) for all $x \in \text{obj}(\mathcal{C})$, the diagrams

$$\begin{array}{ccc}
 1_{\mathcal{D}} \otimes_{\mathcal{D}} F(x) & \xrightarrow{e \otimes \text{id}} & F(1_{\mathcal{C}}) \otimes_{\mathcal{D}} F(x) \\
 \downarrow & & \downarrow \mu \\
 F(x) & \longleftarrow & F(1_{\mathcal{C}} \otimes_{\mathcal{C}} x)
 \end{array}$$

and

$$\begin{array}{ccc}
 F(x) \otimes_{\mathcal{D}} 1_{\mathcal{D}} & \xrightarrow{\text{id} \otimes e} & F(x) \otimes_{\mathcal{D}} F(1_{\mathcal{C}}) \\
 \downarrow & & \downarrow \mu \\
 F(x) & \longleftarrow & F(x \otimes_{\mathcal{C}} 1_{\mathcal{C}})
 \end{array}$$

commute.

Definition 2. A **strong monoidal functor** is a lax monoidal functor $(F, e, \mu_{x,y})$ such that e and $\mu_{x,y}$ are isomorphisms.

Definition 3. Let $(\mathcal{C}, \otimes, 1)$ be a monoidal category. An **algebra object** (A, m, u) in \mathcal{C} is

- an object $A \in \text{obj}(\mathcal{C})$,
- a morphism $m : A \otimes A \rightarrow A$, and
- a morphism $u : 1 \rightarrow A$,

such that

1. (**associativity**) the diagram

$$\begin{array}{ccccc}
 (A \otimes A) \otimes A & \xrightarrow{a} & A \otimes (A \otimes A) & \xrightarrow{\text{id} \otimes m} & A \otimes A \\
 m \otimes \text{id} \downarrow & & & & \downarrow m \\
 A \otimes A & \xrightarrow{\quad m \quad} & & & A
 \end{array}$$

commutes, and

2. (**unitality**) the diagram

$$\begin{array}{ccccc}
 1 \otimes A & \xrightarrow{u \otimes \text{id}} & A \otimes A & \xleftarrow{\text{id} \otimes u} & A \otimes 1 \\
 & \searrow & \downarrow m & \swarrow & \\
 & & A & &
 \end{array}$$

commutes.

Note that if A is an algebra object, and F is a lax monoidal functor, then the lax structure of F translates the multiplication of A into a multiplication on $F(A)$,

$$F(A) \otimes F(A) \xrightarrow{\mu} F(A \otimes A) \xrightarrow{F(m)} F(A).$$

Similarly, there is an induced unit

$$1_{\mathcal{D}} \xrightarrow{e} F(1_{\mathcal{C}}) \xrightarrow{F(u)} F(A).$$

Proposition 4. Let $(F, e, \mu_{x,y}) : (\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}}) \rightarrow (\mathcal{D}, \otimes_{\mathcal{D}}, 1_{\mathcal{D}})$ be a lax monoidal functor. If (A, m, u) is an algebra object in \mathcal{C} , then $(F(A), F(m) \circ \mu_{A,A}, F(u) \circ e)$ is an algebra object in \mathcal{D} .

Proof. For associativity, we need to verify the commutativity of the diagram

$$\begin{array}{ccccc}
 (F(A) \otimes F(A)) \otimes F(A) & \xrightarrow{a} & F(A) \otimes (F(A) \otimes F(A)) & \longrightarrow & F(A) \otimes F(A) \\
 \downarrow & & & & \downarrow \\
 F(A) \otimes F(A) & \xrightarrow{\quad \quad \quad} & & & F(A)
 \end{array}$$

which expands to

$$\begin{array}{ccccc}
(F(A) \otimes F(A)) \otimes F(A) & \xrightarrow{a} & F(A) \otimes (F(A) \otimes F(A)) & \xrightarrow{\text{id} \otimes \mu} & F(A) \otimes F(A \otimes A) & \xrightarrow{\text{id} \otimes F(m)} & F(A) \otimes F(A) \\
\downarrow \mu \otimes \text{id} & & & & \downarrow \mu & & \downarrow \mu \\
F(A \otimes A) \otimes F(A) & \xrightarrow{\mu} & F((A \otimes A) \otimes A) & \xrightarrow{F(a)} & F(A \otimes (A \otimes A)) & \xrightarrow{F(\text{id} \otimes m)} & F(A \otimes A) \\
\downarrow F(m) \otimes \text{id} & & \downarrow F(m \otimes \text{id}) & & & & \downarrow F(m) \\
F(A) \otimes F(A) & \xrightarrow{\mu} & F(A \otimes A) & \xrightarrow{F(m)} & F(A) & & F(A)
\end{array}$$

This diagram is a concatenation of four rectangles. The top left rectangle commutes by associativity of the lax structure on F . The bottom right rectangle is F applied to the associativity diagram of A , so it commutes. The top right and bottom left squares commute by naturality of μ , so the entire diagram commutes.

For the unitality diagram,

$$\begin{array}{ccccc}
1_{\mathcal{D}} \otimes F(A) & \longrightarrow & F(A) \otimes F(A) & \longleftarrow & F(A) \otimes 1_{\mathcal{D}} \\
& & \downarrow & & \\
& & F(A) & &
\end{array}$$

both sides behave the same, so we'll just consider the left triangle, which expands to

$$\begin{array}{ccccc}
1_{\mathcal{D}} \otimes F(A) & \xrightarrow{e \otimes \text{id}} & F(1_{\mathcal{C}}) \otimes F(A) & \xrightarrow{F(u) \otimes \text{id}} & F(A) \otimes F(A) \\
\downarrow & & \downarrow \mu & & \downarrow \mu \\
F(A) & \longleftarrow & F(1_{\mathcal{C}} \otimes A) & \xrightarrow{F(u \otimes \text{id})} & F(A \otimes A) \\
& & \searrow & \swarrow & \\
& & & & F(m)
\end{array}$$

which is a concatenation of two squares atop a (curved) triangle. The left square commutes by unitality of the lax structure of F . The right square commutes by naturality of μ . The bottom triangle is F applied to the unitality diagram of A , so it commutes. Thus, the entire diagram commutes. \square

Definition 5. Let (A, m, u) be an algebra object in a monoidal category $(\mathcal{C}, \otimes, 1)$. A **(left) module** (M, s) **over** A (or A -**module**) in \mathcal{C} is

- an object $M \in \text{obj}(\mathcal{C})$, and
- a morphism $s : A \otimes M \rightarrow M$,

such that

1. (**associativity**) the diagram

$$\begin{array}{ccccc}
(A \otimes A) \otimes M & \xrightarrow{a} & A \otimes (A \otimes M) & \xrightarrow{\text{id} \otimes s} & A \otimes M \\
m \otimes \text{id} \downarrow & & & & \downarrow s \\
A \otimes M & \xrightarrow{\quad\quad\quad s \quad\quad\quad} & & & M
\end{array}$$

commutes, and

2. (**unitality**) the diagram

$$\begin{array}{ccc}
1 \otimes M & \xrightarrow{u \otimes \text{id}} & A \otimes M \\
& \searrow & \swarrow s \\
& & M
\end{array}$$

commutes.

As above, lax structures translate the action of an algebra A on a module M into an action of $F(A)$ on $F(M)$,

$$F(A) \otimes F(M) \xrightarrow{\mu} F(A \otimes M) \xrightarrow{F(s)} F(M).$$

Proposition 6. Let $(F, e, \mu_{x,y}) : (\mathcal{C}, \otimes_{\mathcal{C}}, 1_{\mathcal{C}}) \rightarrow (\mathcal{D}, \otimes_{\mathcal{D}}, 1_{\mathcal{D}})$ be a lax monoidal functor, and let (A, m, u) be an algebra object in \mathcal{C} . If (M, s) is a module over A then $(F(M), F(s) \circ \mu_{A,M})$ is a module over $F(A)$.

Proof. The proof is similar to the proof of Proposition 4, but since I don't know where it's written, I'll write it. For associativity, we need to verify commutativity of the diagram

$$\begin{array}{ccc}
(F(A) \otimes F(A)) \otimes F(M) & \xrightarrow{a} & F(A) \otimes (F(A) \otimes F(M)) \longrightarrow F(A) \otimes F(M) \\
\downarrow & & \downarrow \\
F(A) \otimes F(M) & \xrightarrow{\quad\quad\quad} & F(M)
\end{array}$$

which expands to

$$\begin{array}{ccccccc}
(F(A) \otimes F(A)) \otimes F(M) & \xrightarrow{a} & F(A) \otimes (F(A) \otimes F(M)) & \xrightarrow{\text{id} \otimes \mu} & F(A) \otimes F(A \otimes M) & \xrightarrow{\text{id} \otimes F(s)} & F(A) \otimes F(M) \\
\mu \otimes \text{id} \downarrow & & & & \downarrow \mu & & \downarrow \mu \\
F(A \otimes A) \otimes F(M) & \xrightarrow{\mu} & F((A \otimes A) \otimes M) & \xrightarrow{F(a)} & F(A \otimes (A \otimes M)) & \xrightarrow{F(\text{id} \otimes s)} & F(A \otimes M) \\
F(m) \otimes \text{id} \downarrow & & \downarrow F(m \otimes \text{id}) & & & & \downarrow F(s) \\
F(A) \otimes F(M) & \xrightarrow{\mu} & F(A \otimes M) & \xrightarrow{\quad\quad\quad F(s) \quad\quad\quad} & & & F(M).
\end{array}$$

This diagram is a concatenation of four rectangles. The top left rectangle commutes by associativity of the lax structure on F . The bottom right rectangle is F applied to the associativity diagram of M , so it commutes. The top right and bottom left squares commute by naturality of μ , so the entire diagram commutes.

For unitality, the relevant diagram is

$$\begin{array}{ccc}
1_{\mathcal{D}} \otimes F(M) & \longrightarrow & F(A) \otimes F(M) \\
& \searrow & \swarrow \\
& & F(M)
\end{array}$$

which expands to

$$\begin{array}{ccccc}
1_{\mathcal{D}} \otimes F(M) & \xrightarrow{e \otimes \text{id}} & F(1_{\mathcal{C}}) \otimes F(M) & \xrightarrow{F(u) \otimes \text{id}} & F(A) \otimes F(M) \\
\downarrow & & \downarrow \mu & & \downarrow \mu \\
F(M) & \longleftarrow & F(1_{\mathcal{C}} \otimes M) & \xrightarrow{F(u \otimes \text{id})} & F(A \otimes M) \\
& & \searrow & \swarrow & \\
& & & & F(s)
\end{array}$$

which is a concatenation of two squares atop a (curved) triangle. The left square commutes by unitality of the lax structure of F . The right square commutes by naturality of μ . The bottom triangle is F applied to the unitality diagram of M , so it commutes. Thus, the entire diagram commutes. \square

We've established that lax monoidal functors take algebra objects to algebra objects and modules to modules. Next, we'd like to say something about the *categories* of modules over algebra objects.

Definition 7. Let (A, m, u) be an algebra object in a monoidal category $(\mathcal{C}, \otimes, 1)$, and let (M, s) and (N, t) be modules over A . A morphism $f : M \rightarrow N$ in \mathcal{C} is called a **homomorphism** of A -modules if the diagram

$$\begin{array}{ccc}
A \otimes M & \xrightarrow{s} & M \\
\text{id} \otimes f \downarrow & & \downarrow f \\
A \otimes N & \xrightarrow{t} & N
\end{array}$$

commutes.

Since compositions of module homomorphisms are module homomorphisms, and identity morphisms of modules are module homomorphisms, A -modules and their homomorphisms form a subcategory of \mathcal{C} , denoted ${}_A\text{Mod}$.

Proposition 8. If $f : M \rightarrow N$ is a homomorphism of A -modules and F is a lax monoidal functor, then $F(f)$ is a homomorphism of $F(A)$ -modules.

Proof. We need to verify commutativity of the diagram

$$\begin{array}{ccc}
F(A) \otimes F(M) & \longrightarrow & F(M) \\
\downarrow & & \downarrow \\
F(A) \otimes F(N) & \longrightarrow & F(N)
\end{array}$$

which can be expanded to

$$\begin{array}{ccccc}
F(A) \otimes F(M) & \xrightarrow{\mu} & F(A \otimes M) & \xrightarrow{F(s)} & F(M) \\
\text{id} \otimes F(f) \downarrow & & F(\text{id} \otimes f) \downarrow & & \downarrow F(f) \\
F(A) \otimes F(N) & \xrightarrow{\mu} & F(A \otimes N) & \xrightarrow{F(t)} & F(N).
\end{array}$$

The left square commutes by naturality of μ , and the right square commutes since f is a homomorphism of A -modules and F is a functor. \square

Corollary 9. If $F : \mathcal{C} \rightarrow \mathcal{D}$ is a lax monoidal functor and $A \in \text{obj}(\mathcal{C})$ is an algebra object, then $F({}_A\text{Mod}) \subseteq {}_{F(A)}\text{Mod}$.

We would like to know when this relationship is stronger than an inclusion. The situation is fairly simple for isomorphisms of categories.

Lemma 10. If $(F, e, \mu_{x,y})$ is a strong monoidal functor such that $F : \mathcal{C} \rightarrow \mathcal{D}$ is an isomorphism of categories, then $(F^{-1}, \tilde{e}, \tilde{\mu}_{x,y})$ is a strong monoidal functor, where $\tilde{e} = F^{-1}(e^{-1})$ and $\tilde{\mu}_{x,y} = F^{-1}(\mu_{F^{-1}(x), F^{-1}(y)}^{-1})$.

Proof. First, observe that \tilde{e} and $\tilde{\mu}_{x,y}$ do indeed have the appropriate domains and codomains. The associativity and unitality diagrams can be obtained by inverting the arrows of the corresponding diagram for F and applying F^{-1} to the whole diagram. Thus, they commute. Naturality of $\tilde{\mu}$ follows from similar considerations. \square

Proposition 11. Let $(F, e, \mu_{x,y})$ be a strong monoidal functor such that $F : \mathcal{C} \rightarrow \mathcal{D}$ is an isomorphism of categories, and let $A \in \text{obj}(\mathcal{C})$ be an algebra object. Then F restricts to an isomorphism of categories $F_A : {}_A\text{Mod} \rightarrow {}_{F(A)}\text{Mod}$.

Proof. Since F is an isomorphism, F restricts to an isomorphism ${}_A\text{Mod} \rightarrow F({}_A\text{Mod})$, so by Corollary 9, we just need to show that ${}_{F(A)}\text{Mod} \subset F({}_A\text{Mod})$. Since F^{-1} is also a lax monoidal functor by Lemma 10, we can apply Corollary 9 to F^{-1} and $F(A)$ to obtain the desired inclusion. \square

A similar fact holds for equivalences of categories, but it is more cumbersome.

Lemma 12. Let $(F, e, \mu_{x,y})$ be a strong monoidal functor such that $F : \mathcal{C} \rightarrow \mathcal{D}$ is an equivalence of categories, and let $G : \mathcal{D} \rightarrow \mathcal{C}$ be an inverse equivalence with natural isomorphisms $\alpha : \text{id}_{\mathcal{C}} \Rightarrow GF$ and $\beta : FG \Rightarrow \text{id}_{\mathcal{D}}$. Then $(G^{-1}, \tilde{e}, \tilde{\mu}_{x,y})$ is a strong monoidal functor, where $\tilde{e} = G(e^{-1}) \circ \alpha_{1_{\mathcal{C}}}$ and $\tilde{\mu}_{x,y} = G((\beta_x \otimes \beta_y) \circ \mu_{G(x), G(y)}^{-1}) \circ \alpha_{G(x), G(y)}$.

Proof. Similar to Lemma 10, but with more commutative diagrams to unpack due to the extra natural isomorphisms. \square

Proposition 13. Let $(F, e, \mu_{x,y})$ be a strong monoidal functor such that $F : \mathcal{C} \rightarrow \mathcal{D}$ is an equivalence of categories, and let $A \in \text{obj}(\mathcal{C})$ be an algebra object. Then F restricts to an equivalence of categories $F_A : {}_A\text{Mod} \rightarrow {}_{F(A)}\text{Mod}$.

Proof. Since F is an equivalence, F restricts to an equivalence ${}_A\text{Mod} \rightarrow F({}_A\text{Mod}) \subseteq {}_{F(A)}\text{Mod}$. Similarly, G restricts to an equivalence ${}_{F(A)}\text{Mod} \rightarrow G({}_{F(A)}\text{Mod}) \subseteq {}_{GF(A)}\text{Mod} \cong {}_A\text{Mod}$. Thus, G (together with the natural isomorphisms α, β from Lemma 12) provides an inverse equivalence for F_A . \square